

Web Infrastructures

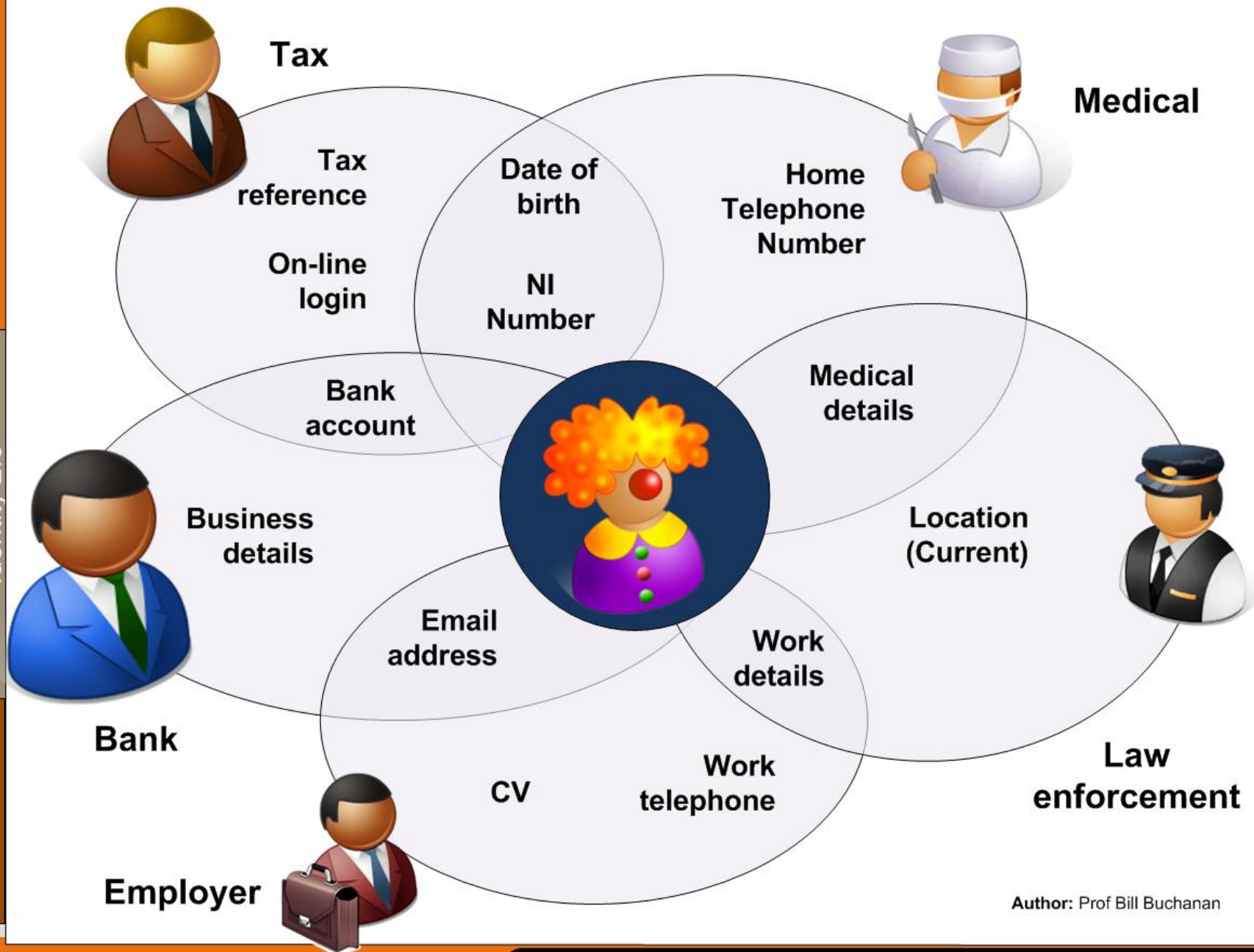
- Provide an overview of Web-based architectures, especially in authentication and access control.
- Define key protocols involved in next generation Web-based infrastructures, such as Kerberos and SOAP over HTTP.
- Define scalable authentication infrastructures and protocols.
- Investigate scaleable and extensible architectures, including using LDAP.



Web Infrastructure



Identity 2.0



Author: Prof Bill Buchanan

Web Infrastructure



SOAP over HTTP

```
<soap:Envelope
  xmlns:xsi="http://www.w3.org/2001/XMLSchema-instance"
  xmlns:xsd="http://www.w3.org/2001/XMLSchema" xmlns:soap="http://
schemas.xmlsoap.org/soap/envelope/">
  <soap:Body>
    <CalcRootResponse xmlns="http://MyMath.com/maths">
      <CalcRootResult>9</CalcRootResult>
    </CalcRootResponse>
  </soap:Body>
</soap:Envelope>
```

**Client
(Windows)**



**SOAP supports the encapsulation of
messages and objects between
different system types**



**Server
(Windows)
- Web Service**



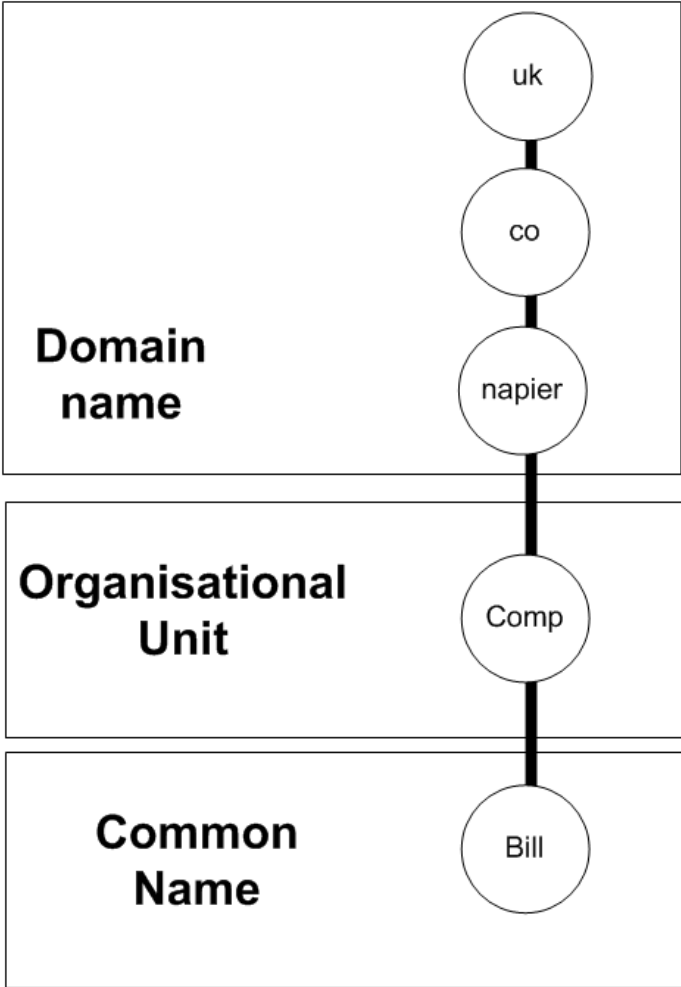
Server (Linux)

```
<double xmlns="http://MyMath.com/maths">3</double>
```


Web Infrastructure



LDAP



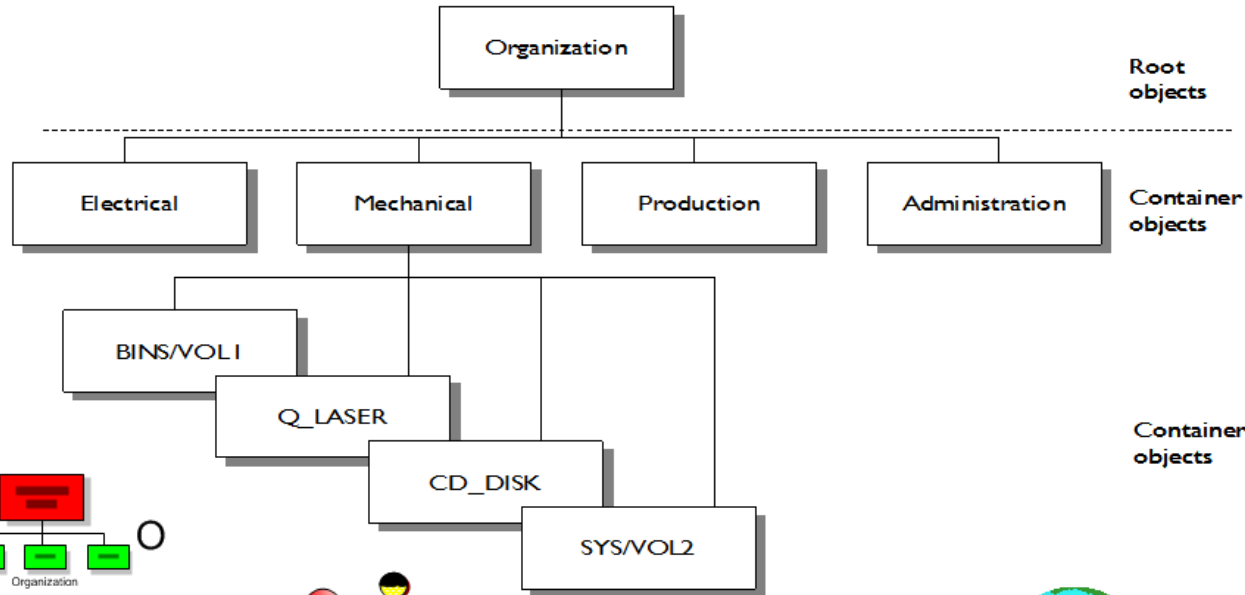
dn: dc=napier,dc=ac,dc=uk
ou: Comp
cn: Bill

Access to Fred's folder
Identifier for Fred login
Identifier for Fred

cn=Fred Folder,ou=people,dc=fake,dc=com
uid=fred,ou=people,dc=fake,dc=com
cn=fred,ou=people,dc=fake,dc=com

Author: Prof Bill Buchanan

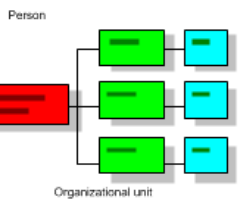
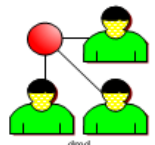
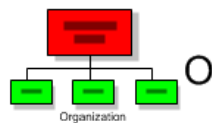
`ldap://ldap.example.com/cn=Bill,dc=napier,dc=ac,dc=uk`



Root objects

Container objects

Container objects



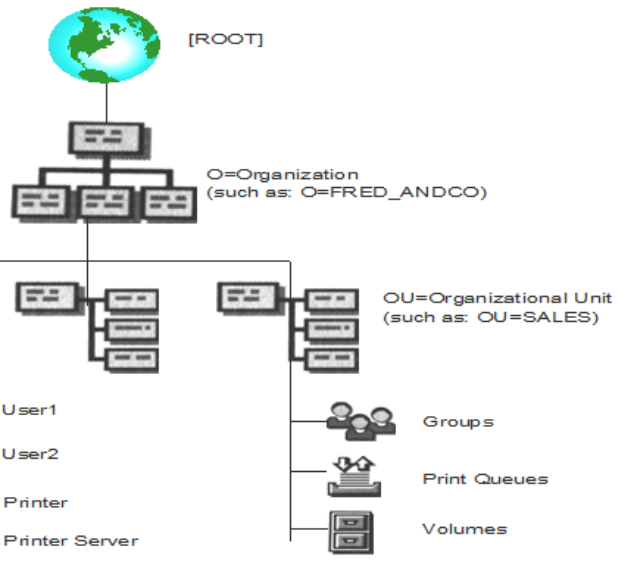
OU



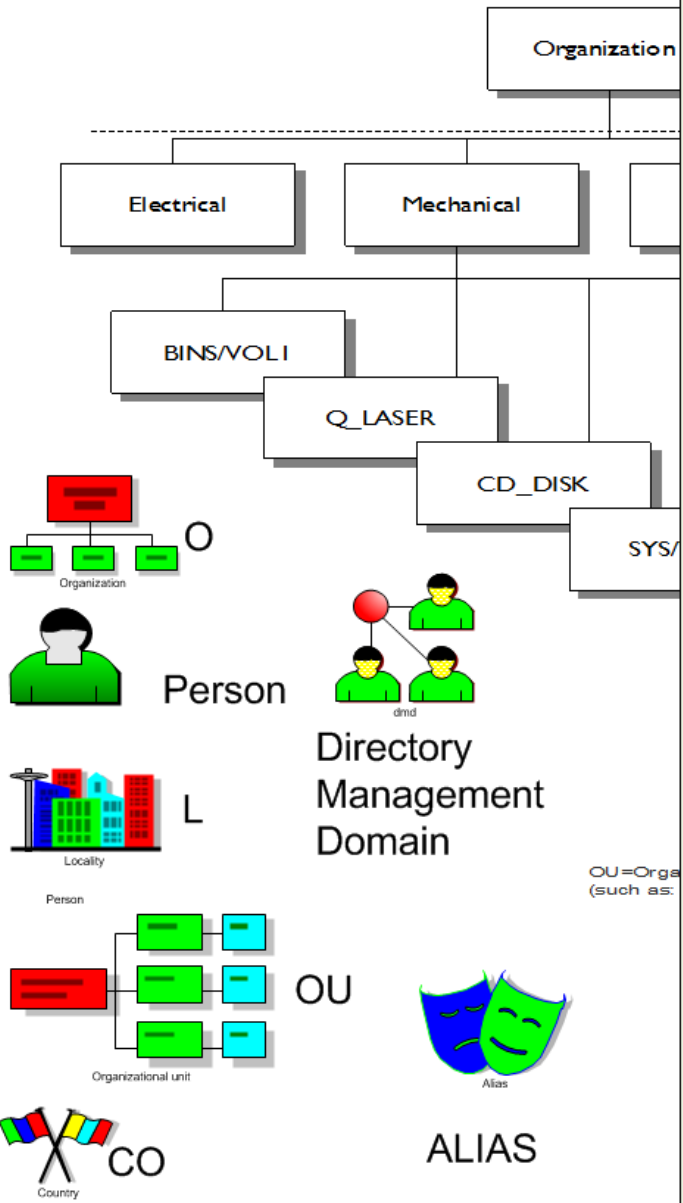
ALIAS



CO



Author: Prof Bill Buchanan

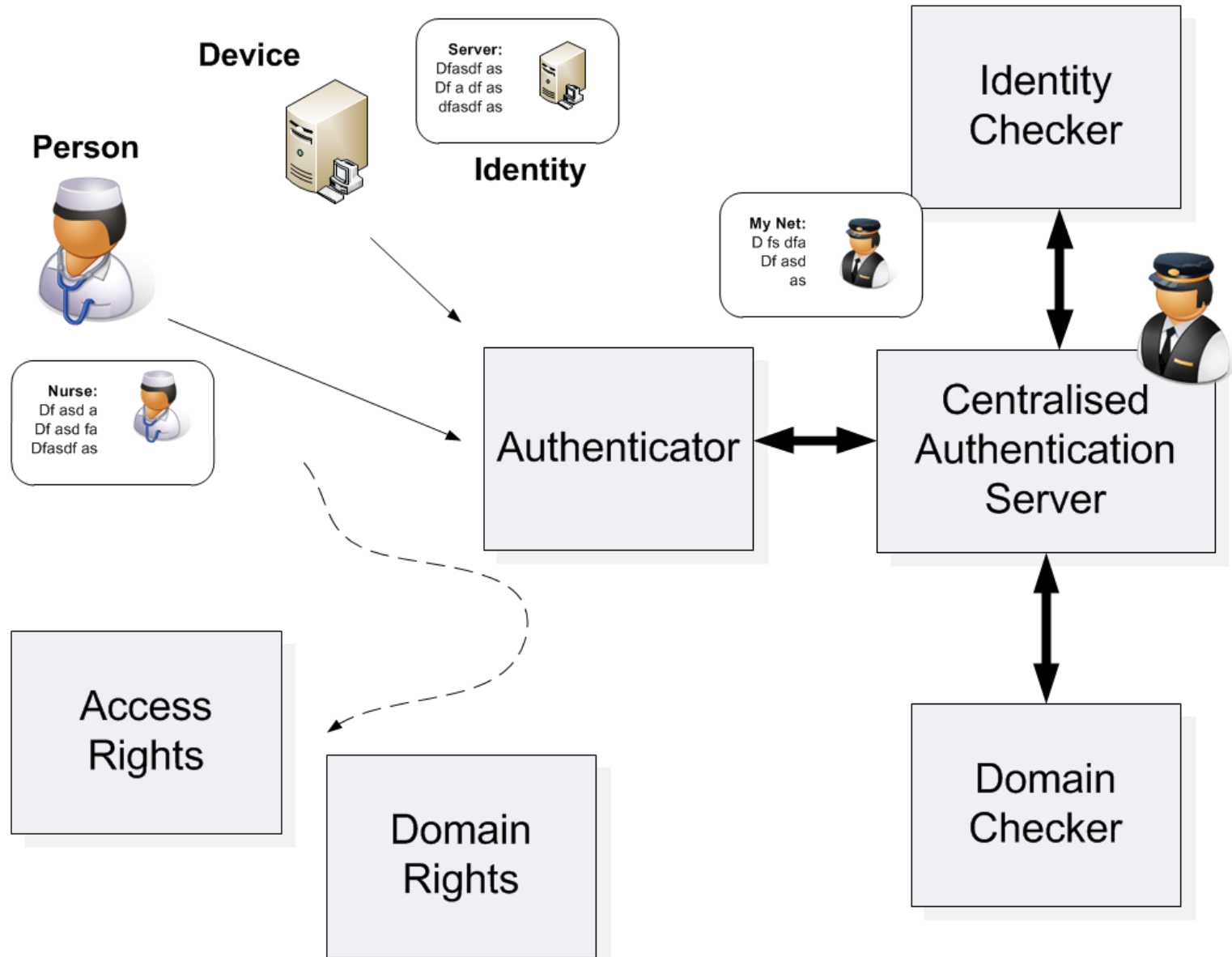


```
dn: ou=people,dc=fake,dc=com
objectClass: organizationalUnit
ou: people
```

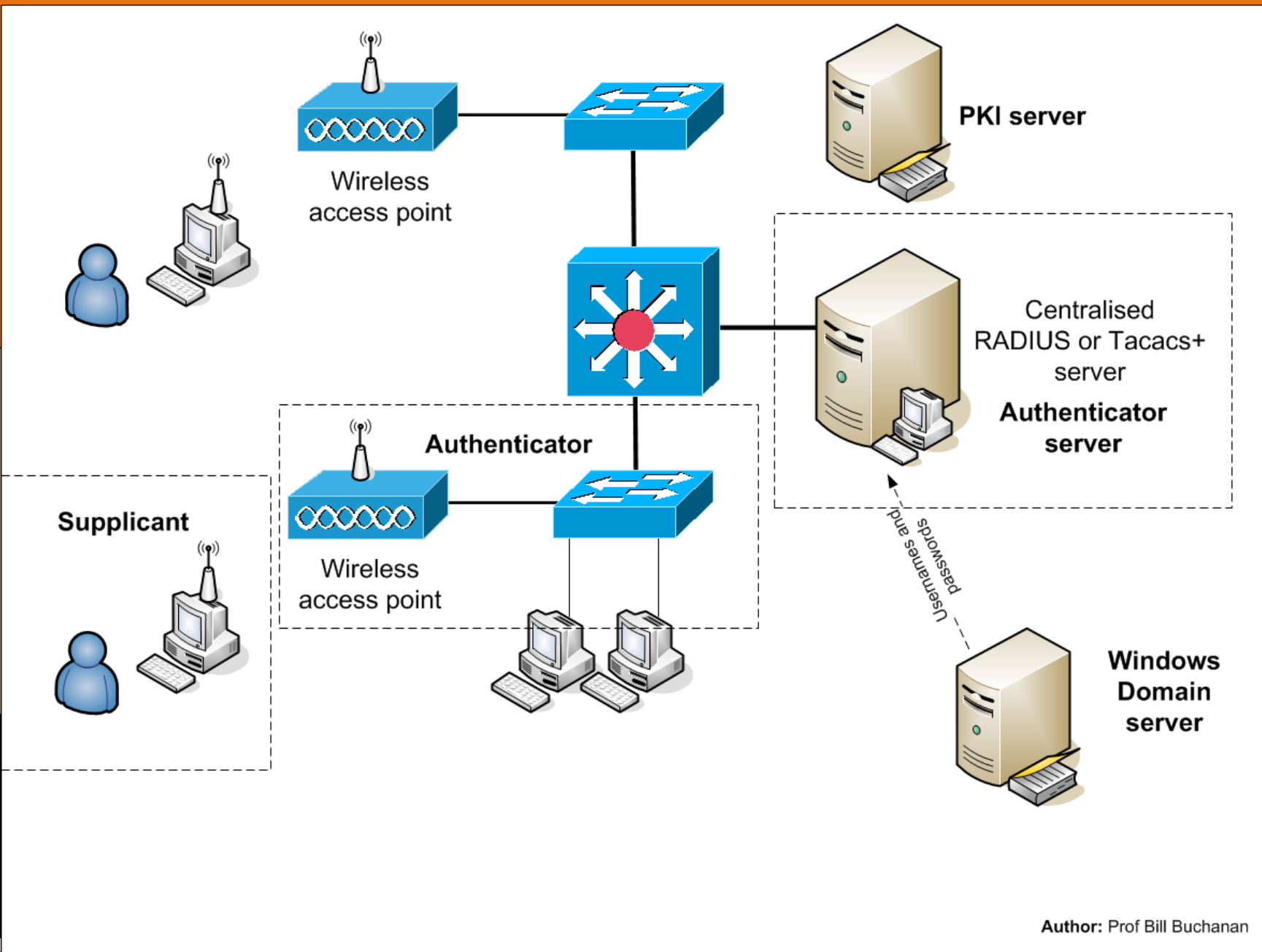
```
dn: ou=groups,dc=fake,dc=com
objectClass: organizationalUnit
ou: groups
```

```
dn: uid=fred, ou= people, dc=fake, dc=com
objectClass: inetOrgPerson
objectClass: posixAccount
objectClass: shadowAccount
uid: fred
givenname: Fred
sn: Fredaldo
cn: Freddy Fredaldo
telephonenumber: 45511332
roomnumber: C.63
o: Fake Inc
mailRoutingAddress: f.smith@fake.com
mailhost: smtp.fake.com
userpassword: {crypt}ggHi99x
uidnumber: 5555
gidnumber: 4321
homedirectory: /user/fred
loginshell: /usr/local/bin/bash
```

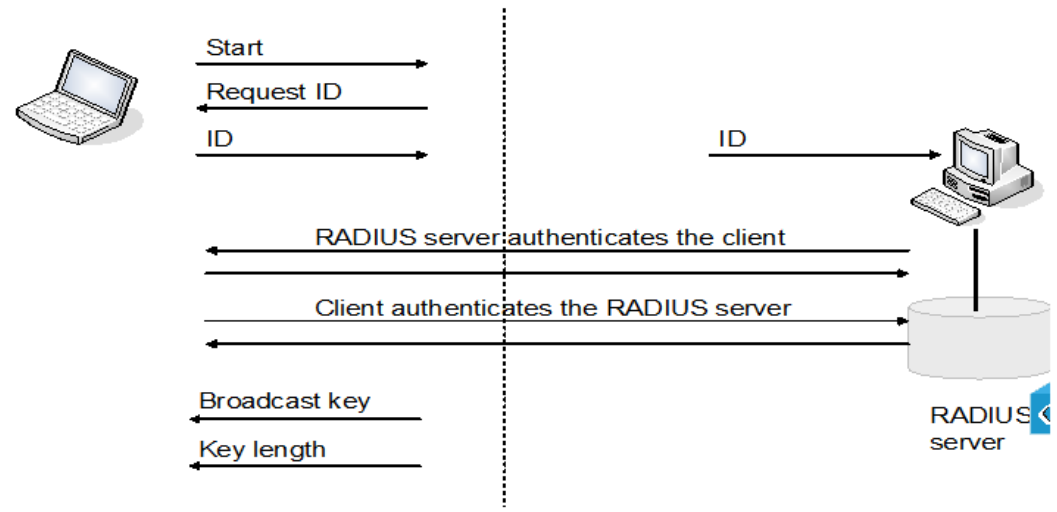
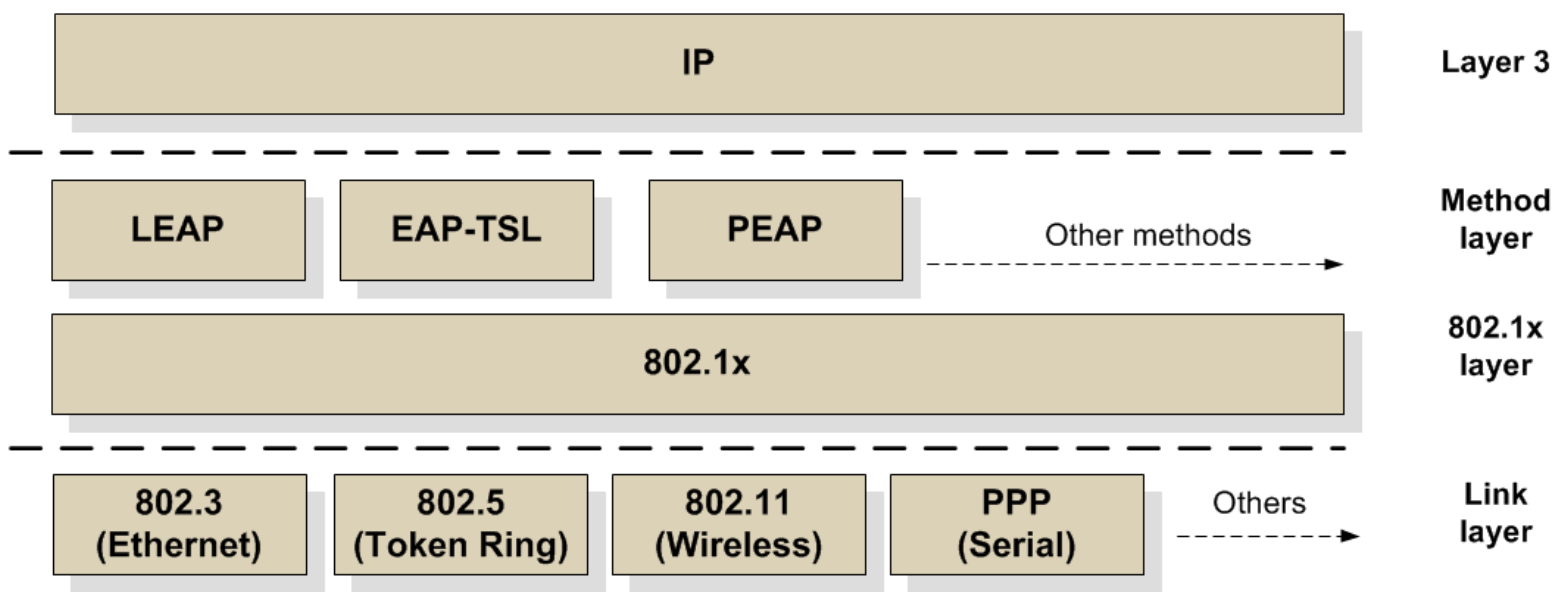
```
dn: cn=example,ou=groups, dc=fake,dc=com
objectClass: posixGroup
cn: example
gidNumber: 10000
```

Author: Prof Bill Buchanan

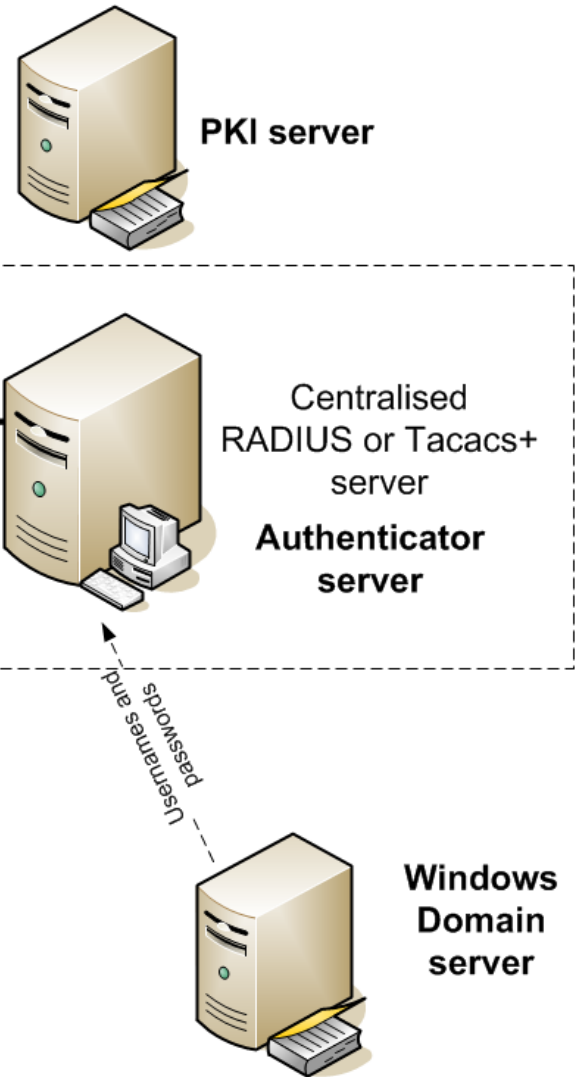


Author: Prof Bill Buchanan



(a)

- **Brute-forcing of user credentials.** A malicious user can continually access the RADIUS server with a range of user ID and associated passwords, and RADIUS may eventually return a success authentication if a match is found.
- **Denial of service.** RADIUS uses UDP, which is connectionless, thus it is difficult to determine malicious from non-malicious UDP packets on ports 1812 and 1813.
- **Session replay.** There is very little authentication of the messages involved in RADIUS, thus malicious users can reply valid ones back into the next at future times.
- **Spoofed packet injection.** There is very little authentication of data packets built into RADIUS, and it can thus suffer from spoofed packet injection.
- **Response Authenticator Attack.** RADIUS uses an MD5-based hash for the Response Authenticator, thus if an intruder captures a valid Access-Request, Access-Accept, or Access-Reject packet sequence, they can launch a brute force attack on the shared secret. This is because the intruder can compute the MD5 hash for (Code+ID+Length+RequestAuth+Attributes), as most of the parts of the Authenticator are known, and can thus focus on the shared secret key.
- **Password Attribute-Based Shared Secret Attack.** Intruders can determine the share secret key but attempting to authenticate using a known password and then capturing the resulting Access-Request packet. After this they can then XOR the protected portion of the User-Password attribute with the password that they have used. A brute-force attack can then be done on the shared secret key
- **Shared Secret.** The basis methodology of RADIUS is that the same shared secret by many clients. Thus weakly protected clients could reveal the secret key.



Author: Prof Bill Buchanan

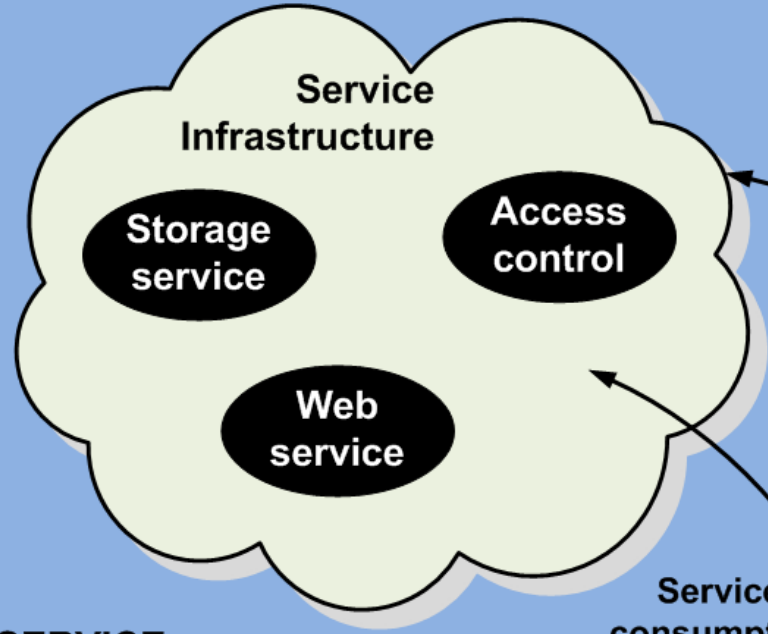
Web Infrastructure



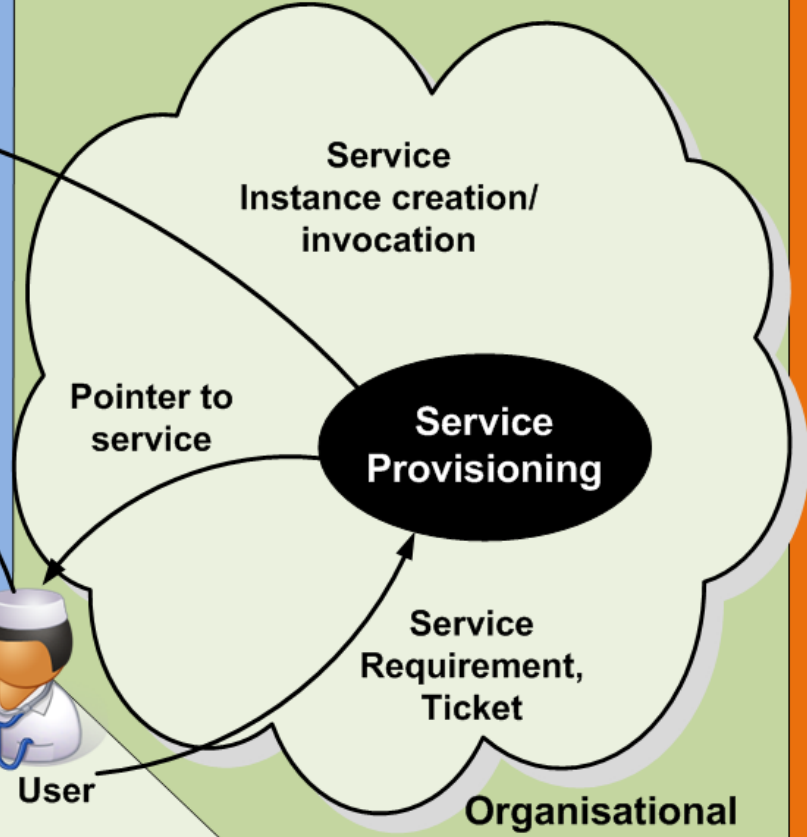
Federated Identity
Management

FIM

Introduction



SERVICE PROVISION

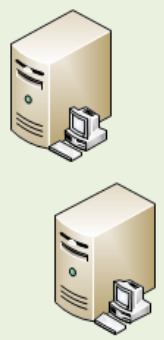


Organisational Infrastructure

SERVICE RIGHTS



Service consumption



Identity credentials

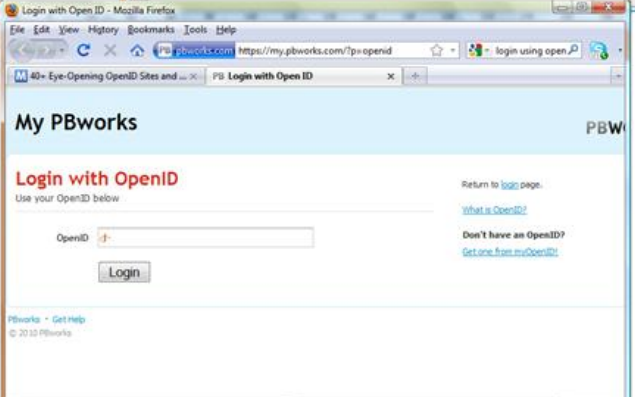
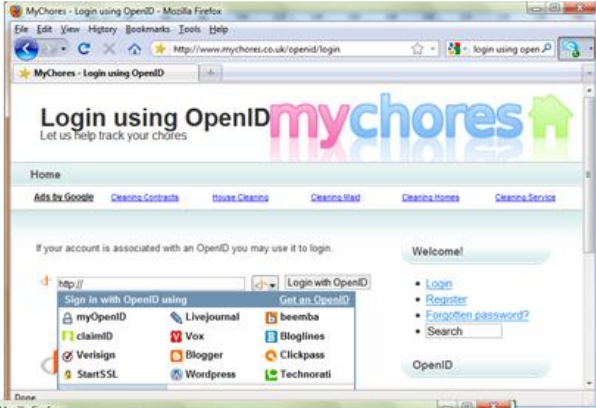
Ticket

Federated Identity Management



Prof Bill Buchanan

Next-generation Web infrastructure



Authenticated into the OpenID infrastructure



<http://billbuchanan.myopenid.com/>

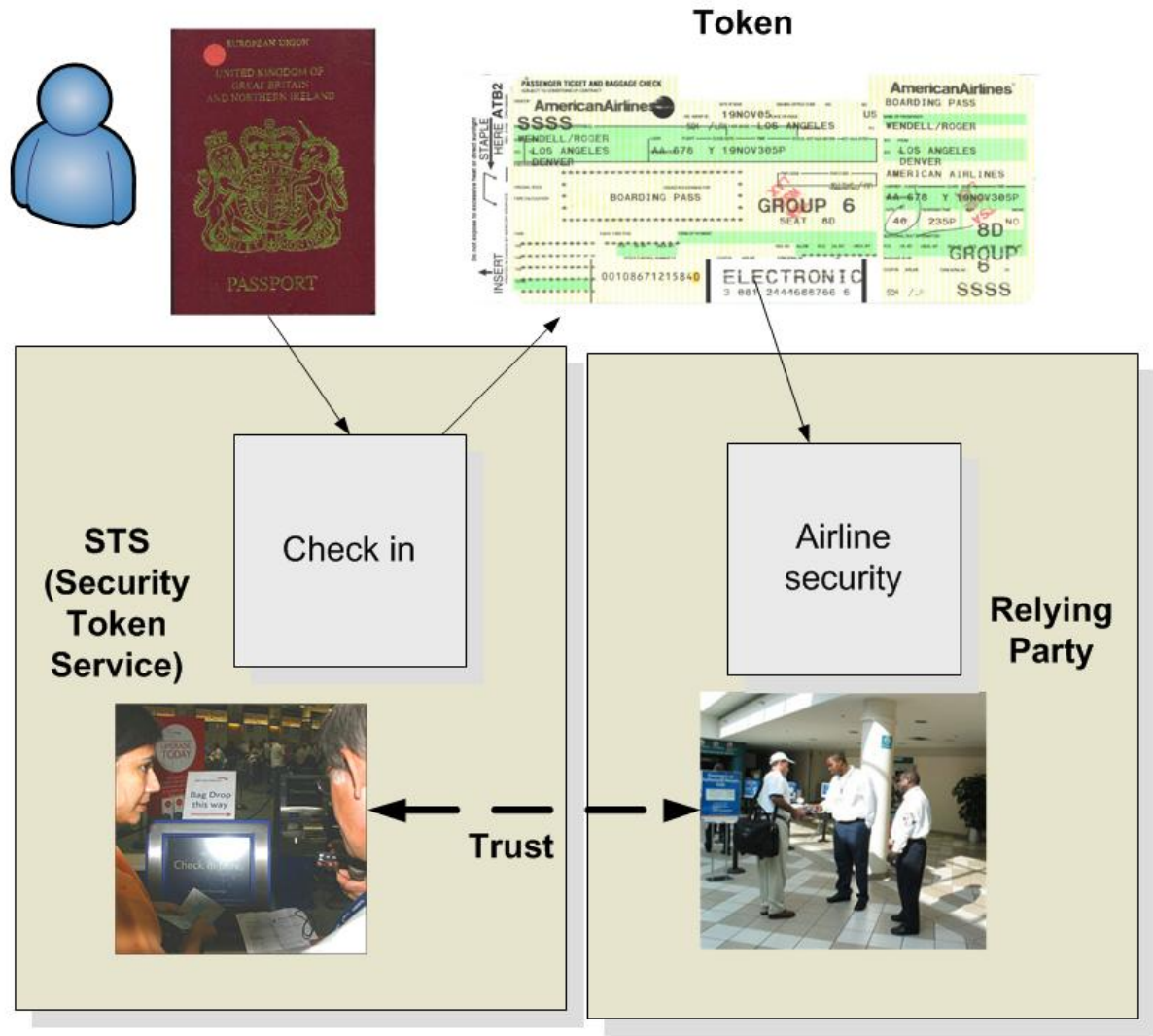


OpenID

Web Infra

Others:

- AOL – openid.aol.com/screenname
- Flickr (Flickr) – www.flickr.com/photos/username
- LiveDoor profile.livedoor.com/username
- Orange (France Telecom) – <http://openid.orange.fr>
- Yahoo (Yahoo!) – <http://openid.yahoo.com>
- WordPress.com – username.wordpress.com



Token

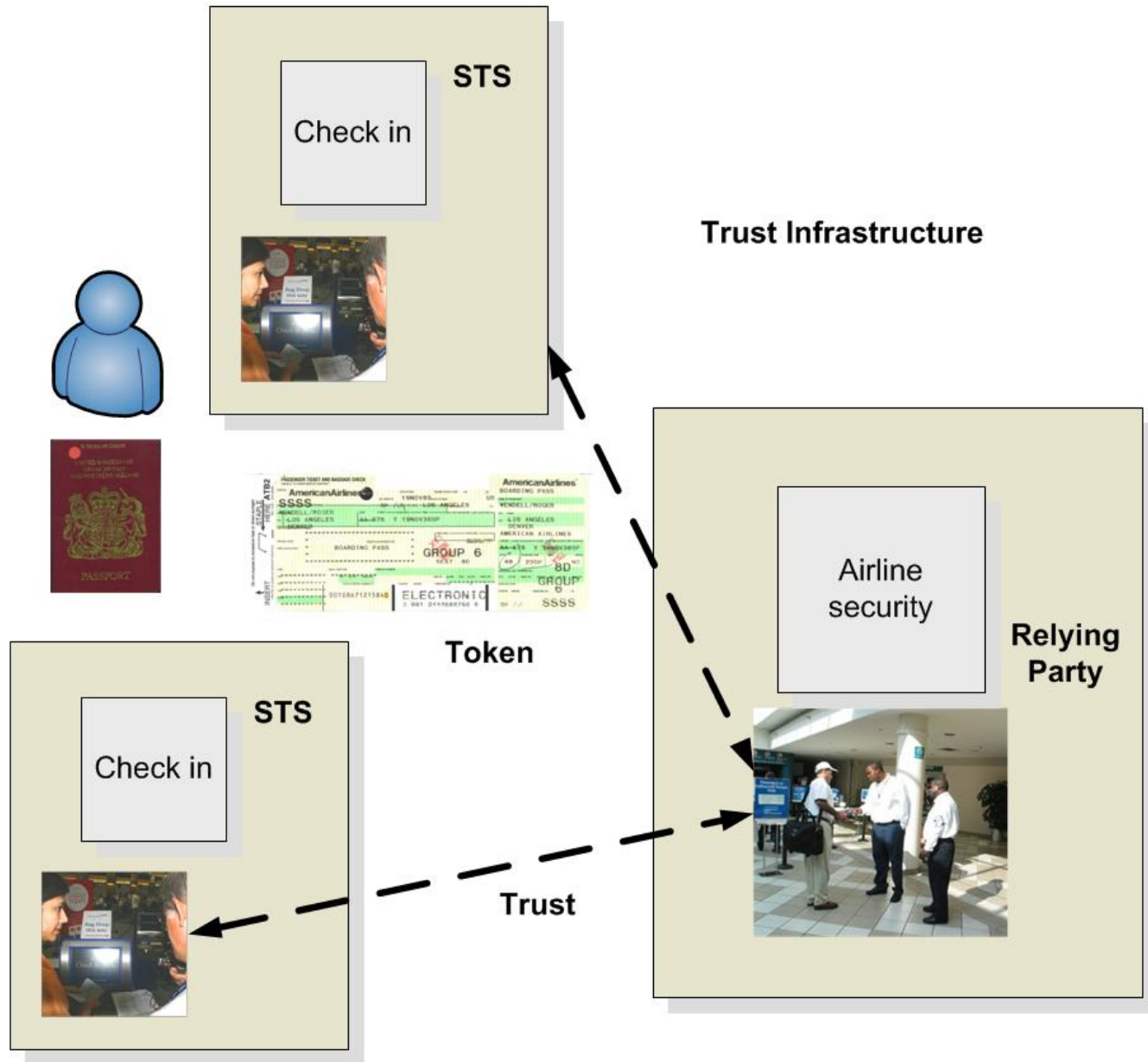
STS
(Security
Token
Service)

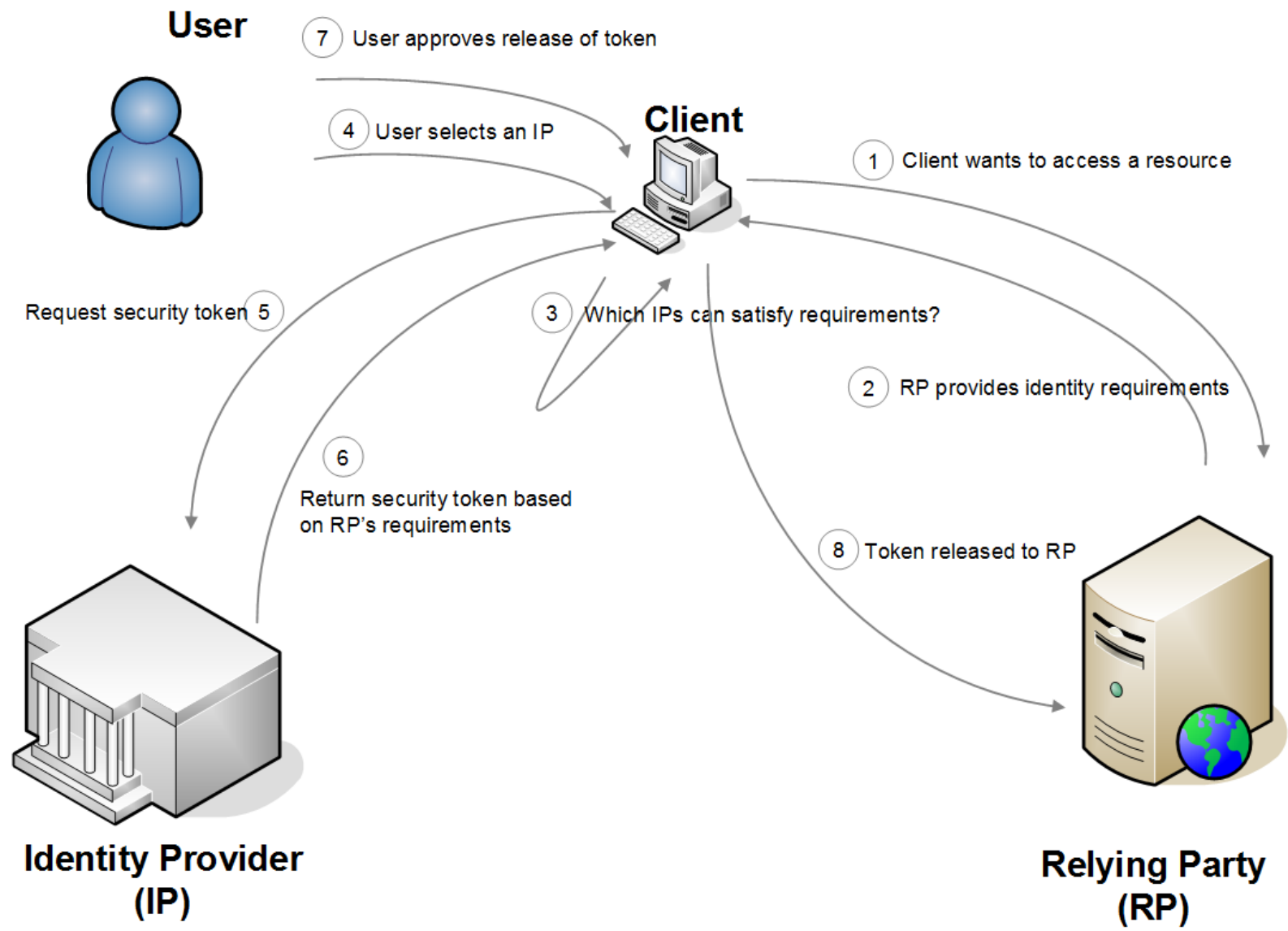
Check in

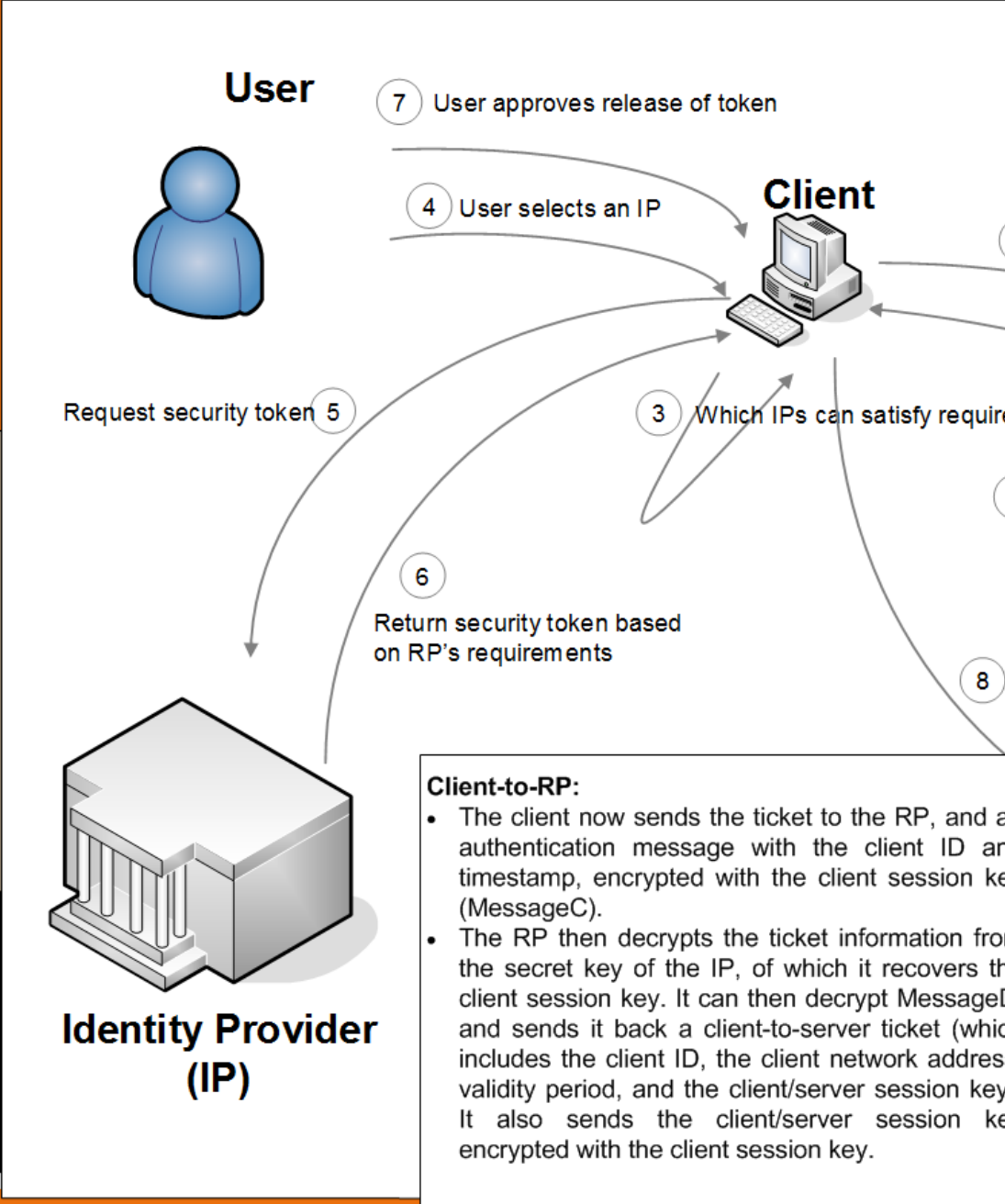
Airline
security

Relying
Party

Trust







Client to IP:

- A user enters a username and password on the client.
- The client performs a one-way function on the entered password, and this becomes the secret key of the client.
- The client sends a cleartext message to the IP requesting services on behalf of the user.
- The IP checks to see if the client is in its database. If it is, the IP sends back a session key encrypted using the secret key of the user (MessageA). It also sends back a ticket which includes the client ID, client network address, ticket validity period, and the client/TGS (Ticket Granting Server) session key encrypted using the secret key of the IP (MessageB).
- Once the client receives messages A and B, it decrypts message A to obtain the client/TGS session key. This session key is used for further communications with IP.

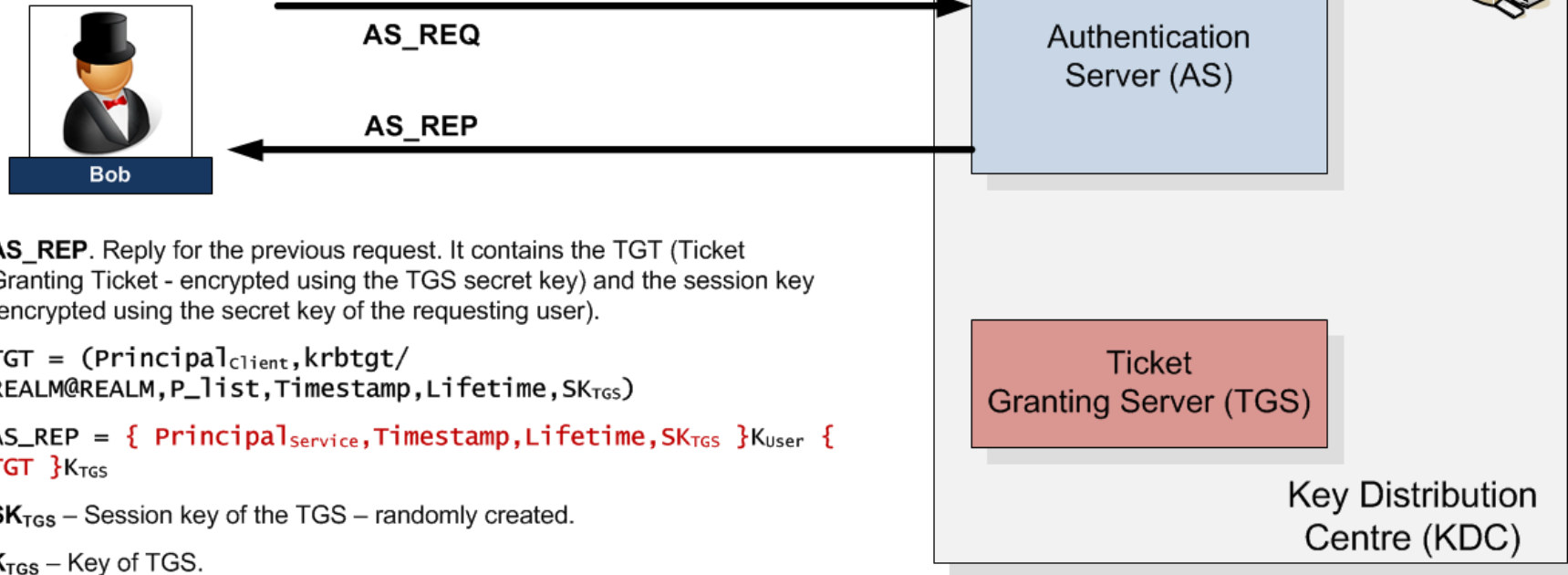
Client-to-RP:

- The client now sends the ticket to the RP, and an authentication message with the client ID and timestamp, encrypted with the client session key (MessageC).
- The RP then decrypts the ticket information from the secret key of the IP, of which it recovers the client session key. It can then decrypt MessageD, and sends it back a client-to-server ticket (which includes the client ID, the client network address, validity period, and the client/server session key). It also sends the client/server session key encrypted with the client session key.

AS_REQ is the initial user authentication request. This message is directed to the KDC component known as Authentication Server (AS).

AS_REQ = (
 Principal_{Client}, **Principal**_{Service}, **IP_list**, **Lifetime**)

Eg **Principal**_{Client} = Principal for user (such as fred@home.com), **IP_list** = all IP address which will use the ticket (may be null if behind NAT), **lifetime** = require life of the ticket.



AS_REP. Reply for the previous request. It contains the TGT (Ticket Granting Ticket - encrypted using the TGS secret key) and the session key (encrypted using the secret key of the requesting user).

TGT = (**Principal**_{Client}, **krbtgt/REALM@REALM**, **P_list**, **Timestamp**, **Lifetime**, **SK_{TGS}**)

AS_REP = { **Principal**_{Service}, **Timestamp**, **Lifetime**, **SK_{TGS}** }**K_{User}** { **TGT** }**K_{TGS}**

SK_{TGS} – Session key of the TGS – randomly created.

K_{TGS} – Key of TGS.

K_{User} – Secret key of Bob.

Note:

{ **Message** } – The curly brackets identify an encrypted message.

(**Message**) – The round brackets identify a non-encrypted message.



$AS_REQ = (Principal_{client}, Principal_{service}, IP_list, Lifetime)$

AS_REQ

$TGT = (Principal_{client}, krbtgt/REALM@REALM, P_list, Timestamp, Lifetime, SK_{TGS})$

$AS_REP = \{ Principal_{service}, Timestamp, Lifetime, SK_{TGS} \}_{K_{user}} \{ TGT \}_{K_{TGS}}$

AS_REP

Bob now needs
a service ticket

TGS_REQ

$Authenticator = \{ Principal_{client}, Timestamp \}_{SK_{TGS}}$

$TGS_REQ = (Principal_{service}, Lifetime, Authenticator) \{ TGT \}_{K_{TGS}}$

TGS_REP

$T_{service} = (Principal_{client}, Principal_{service}, IP_list, Timestamp, Lifetime, SK_{service})$

$TGS_REP = \{ Principal_{service}, Timestamp, Lifetime, SK_{service} \}_{SK_{TGS}} \{ T_{service} \}_{K_{service}}$



Bob

SK_{TGS} – Session key of the TGS – randomly created.

K_{TGS} – Key of TGS.

K_{user} – Secret key of Bob.

$SK_{service}$ – Secret key of the service

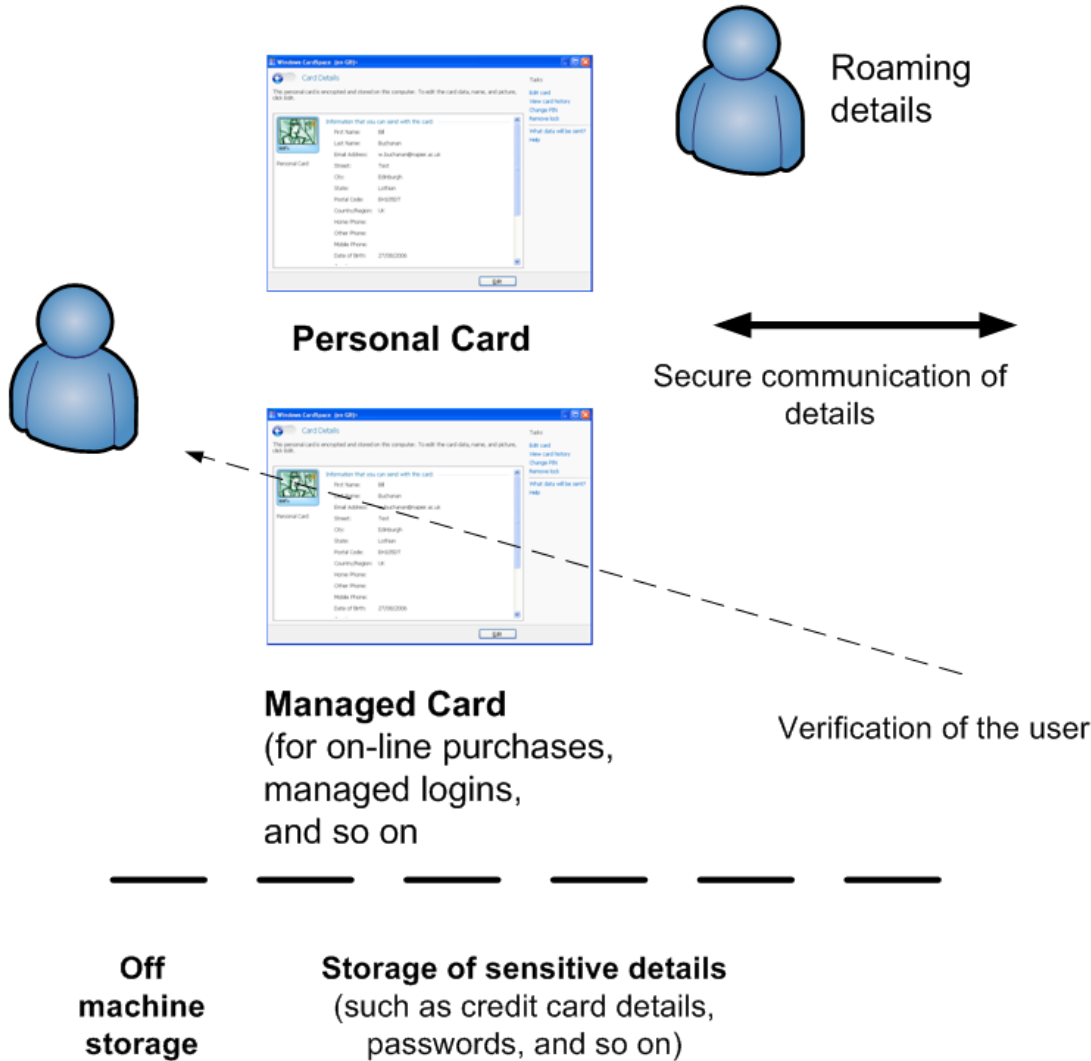
Authentication
Server (AS)

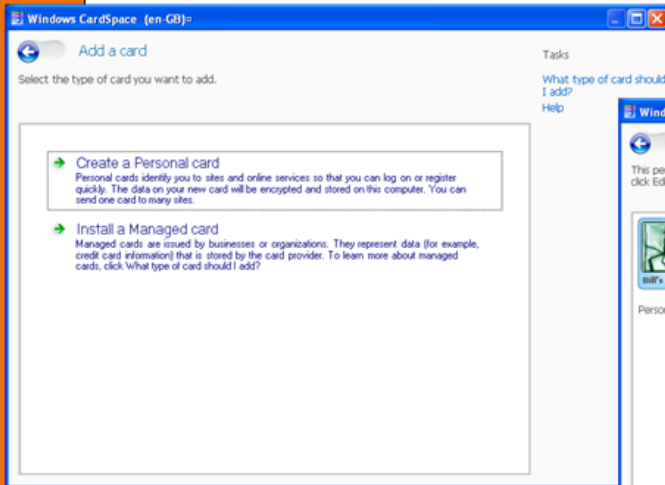
Ticket
Granting Server (TGS)

Key Distribution
Centre (KDC)



Author: Prof Bill Buchanan





Personal cards:

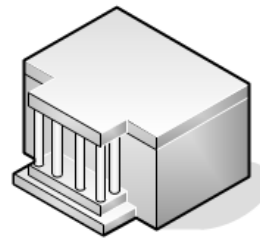
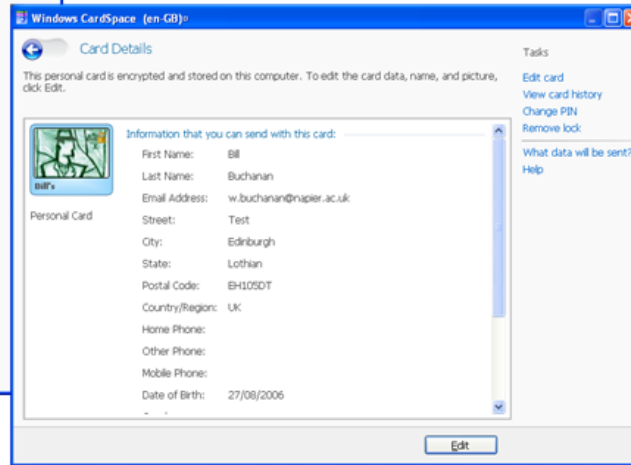
- Created by the person.
- Encrypted.

Personal information:

Name, addresses, phone numbers, date of birth, and gender.

Additional:

Card name, card picture, and card creation date and a history of the sites where this card was used.



Managed Cards:

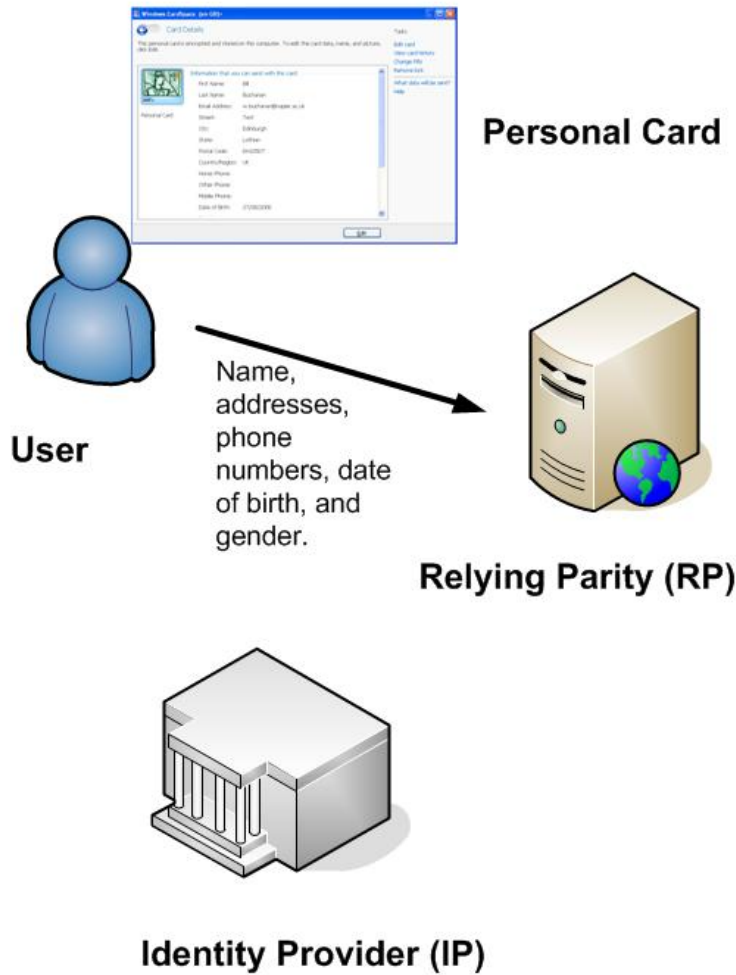
- Created by identity provider.
- Encrypted.

Information:

Maintained by IP that provides card. Stored at site.

Some info on local machine (Card name, when installed, Valid until date, History of card)

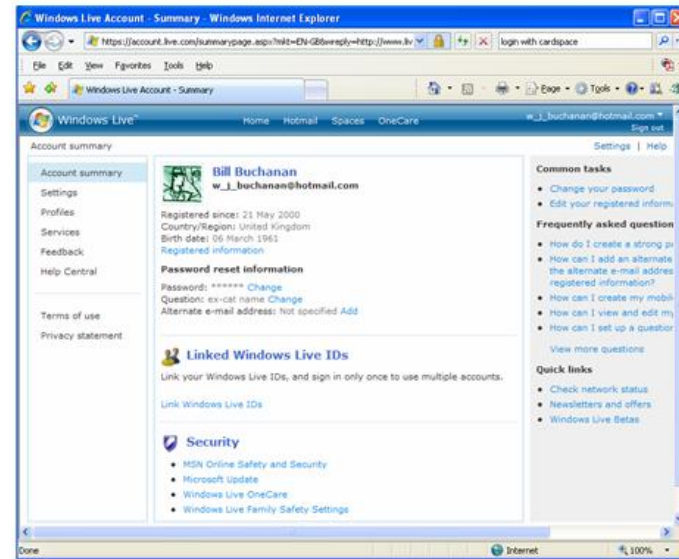
Author: Prof Bill Buchanan



Personal Card

Relying Parity (RP)

Identity Provider (IP)



Author: Prof Bill Buchanan


Windows CardSpace (en-GB)

Select a card to preview

Tasks

To see or edit card data before you send it, select a card, and then click Preview. To create a new card, click Add a card.

Your cards:



Bill's

Add a card

Windows CardSpace (en-GB)

Card Details


This personal card is encrypted and stored on this computer. To edit the card data, name, and picture, click Edit.

Tasks

- Edit card
- View card history
- Change PIN
- Remove lock

What data will be sent?
Help

Information that you can send with this card:



Bill's

Personal Card

First Name: Bill
Last Name: Buchanan
Email Address: w.buchanan@napier.ac.uk
Street: Test
Edinburgh
Lothian
EH105DT

Manage your Information Card - Windows Internet Explorer

https://login.live.com/beta/managecards.srf?wa=wsignin1.0&wreply=http://www...

File Edit View Favorites Tools Help

Manage your Information Card

Windows Live

Sign up

Help

Windows Live ID

Works with MSN, Office Live, and Microsoft Passport sites

Manage your Information Card (Beta)

Windows Live ID now supports Information Cards. By replacing pass... Information Card can help prevent phishing and identity theft.

To manage your Information Card, sign in with your password.

Sign in Cancel

Learn more about Information Cards.

©2007 Microsoft Corporation About Privacy Trademarks Account Help

Windows CardSpace (en-GB)


Choose a card to send to: Microsoft Corporation

To see or edit card data before you send it, select a card, and then click Preview. To create a new card, click Add a card.

Tasks

- Duplicate card
- Delete card
- Add a card
- Back up cards
- Restore cards
- Preferences
- Delete all cards
- Disable Windows CardSpace
- Which card should I send?
- Help
- Learn more about this site

Cards you've sent to this site:



Bill's

Your other cards:

Add a card

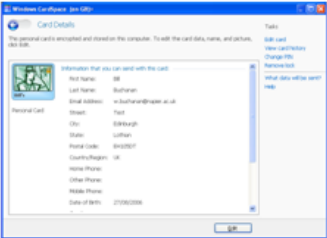
Preview

Send

Web Infrastructure



WS-*



Identity

User


Identity selector



SAML (Security Assertion Markup Language) Or Custom

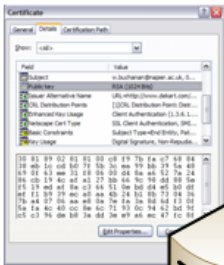
WS-Security Policy

WS-Security




Relying Parity (RP)

Security Token Service (STS)



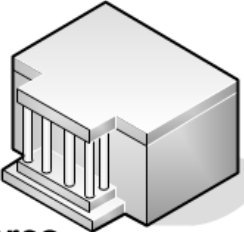
Identity Provider (IP)



X509 Certificate

PKI Server

Digital Certificate Granter (Verisign)



Kerberos

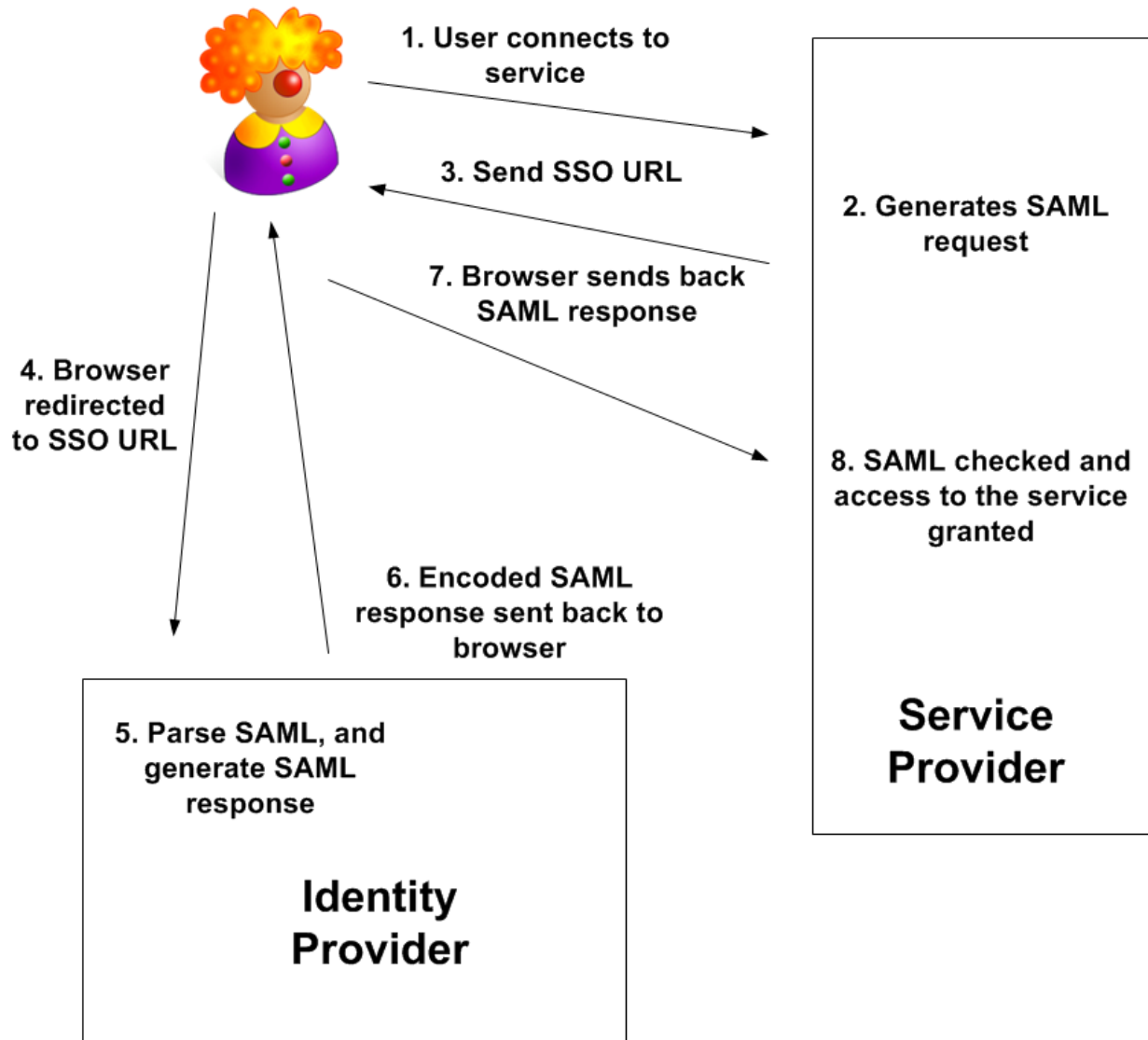
Open XML standards:

WS-*

WS-Trust, WS-Metadata Exchange Framework

```
1 <Assertion ID="_a75adf55-01d7-40cc-929f-dbd8372ebdfc"
2   IssueInstant="2003-04-17T00:46:02Z" Version="2.0"
3   xmlns="urn:oasis:names:tc:SAML:2.0:assertion">
4   <Issuer>
5     example.com
6   </Issuer>
7   <Subject>
8     <NameID
9       Format=
10        "urn:oasis:names:tc:SAML:1.1:nameid-format:emailAddress">
11       Alice@example.com
12     </NameID>
13     <SubjectConfirmation
14       Method="urn:oasis:names:tc:SAML:2.0:cm:sender-vouches"/>
15   </Subject>
16   <Conditions NotBefore="2003-04-17T00:46:02Z"
17     NotOnOrAfter="2003-04-17T00:51:02Z">
18     <AudienceRestriction>
19       <Audience>
20         example2.com
21       </Audience>
22     </AudienceRestriction>
23   </Conditions>
24   <AttributeStatement>
25     <saml:Attribute
26       xmlns:x500=
27        "urn:oasis:names:tc:SAML:2.0:profiles:attribute:X500"
28       NameFormat=
29        "urn:oasis:names:tc:SAML:2.0:attrname-format:uri"
30       Name="urn:oid:2.5.4.20"
31       FriendlyName="telephoneNumber">
32       <saml:AttributeValue xsi:type="xs:string">
33         +1-888-555-1212
34       </saml:AttributeValue>
35     </saml:Attribute>
36   </AttributeStatement>
37 </Assertion>
```

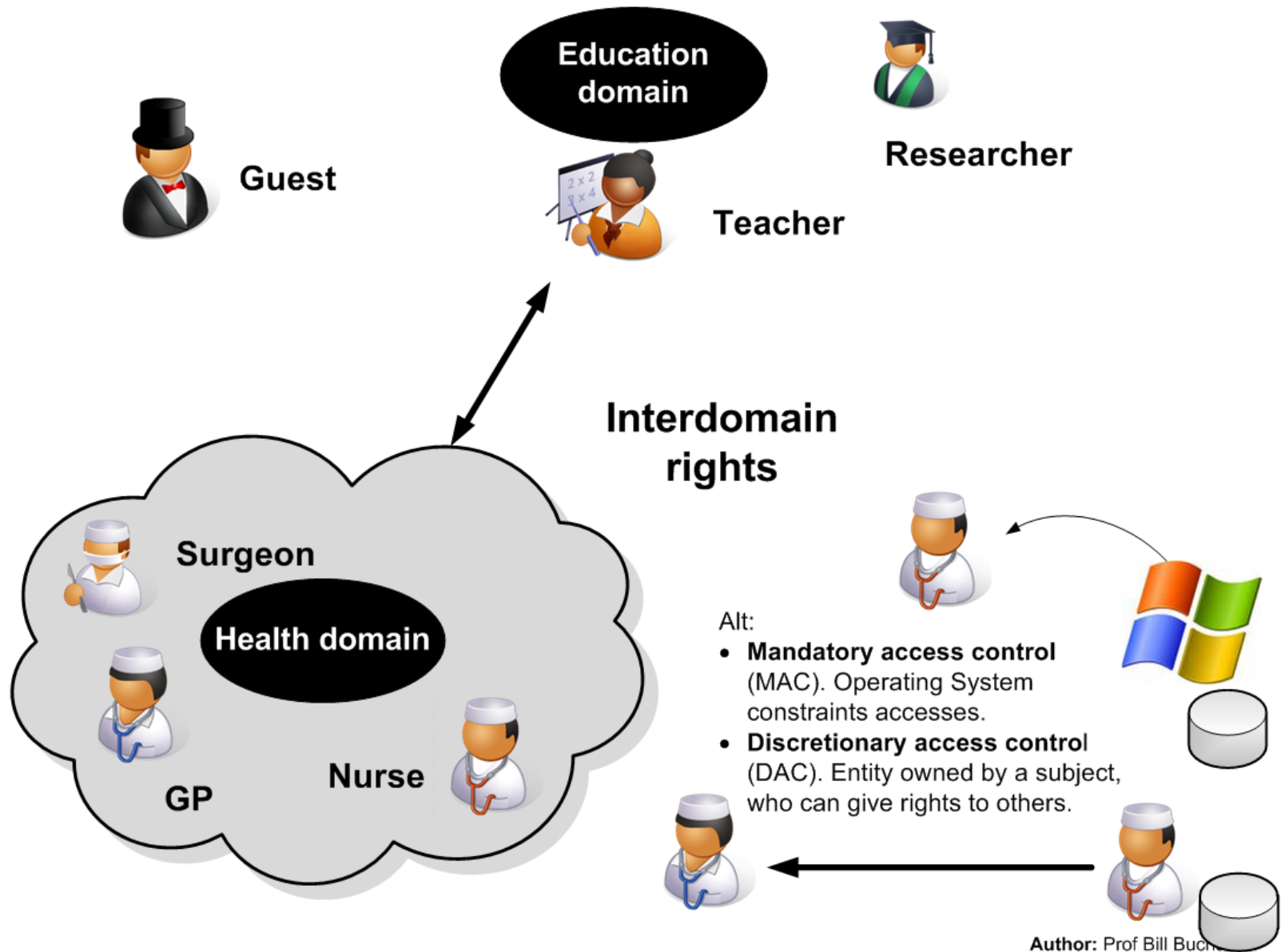
- **SAML Assertions.** These are: **Authentication** assertions (which assert that the user have proven their identity); **Attribute** assertions (which contains information about the user, such as when their limits are); and **Authorization** decision assertions (these define when the user can actually do).
- **Protocol.** This defines method that SAML uses to get gets assertions, such as using SOAP over HTTP (which is the most common method at the present).
- **Binding.** This defines how SAML message are exchanged, such as with SOAP messages.



Web Infrastructure



Access Control



Author: Prof Bill Buch

Web Infrastructures

- Provide an overview of Web-based architectures, especially in authentication and access control.
- Define key protocols involved in next generation Web-based infrastructures, such as Kerberos and SOAP over HTTP.
- Define scalable authentication infrastructures and protocols.
- Investigate scaleable and extensible architectures, including using LDAP.

